

Nonexistence of nontrivial tight 8-design

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Outline

- ▶ t -design
- ▶ Fisher type lower bound
- ▶ Tight $2s$ -design

t -design

Denote set $\{1, \dots, v\}$ by $[v]$.

\mathcal{B} is a t - (v, k, λ) design, t -design in short, if

- ▶ \mathcal{B} is a subset of $\binom{[v]}{k}$,
- ▶ for any $S \in \binom{[v]}{t}$, $\#\{B \in \mathcal{B} : S \subseteq B\} = \lambda > 0$.

Elements of \mathcal{B} are called **blocks**.

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$\mathcal{B} = \binom{[v]}{k}$ is a **trivial** design.

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- ▶ $t = 2$. Fano plane.

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A **Steiner system** is a t -design with $\lambda = 1$ and $t \geq 2$.

Parameter

Let \mathcal{B} be a t -(v, k, λ) design. Then for any subsets $I, J \subseteq [v]$ with $I \cap J = \emptyset$, $|I| + |J| \leq t$,

$$\#\{B \in \mathcal{B} : I \subset B, J \subset \overline{B}\} = \lambda \frac{\binom{v-|I|-|J|}{k-|I|}}{\binom{v-t}{k-t}}.$$

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- ▶ \mathcal{B} is a $(t-1)$ -design.
- ▶ the complementary design $\tilde{\mathcal{B}} = \{\overline{B} : B \in \mathcal{B}\}$ is a t -design.

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Dijen K. Ray-Chaudhuri & Richard M. Wilson. $|\mathcal{B}| \geq \binom{v}{s}$ for $2s$ -designs.

$$\mathbf{M} \in \mathbf{F}_2^{\binom{[v]}{s} \times \mathcal{B}}$$

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\mathbf{M} has full row rank.

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Assume $v \geq 2k$.

Classification of tight $2s$ -designs

- ▶ $s = 2$. Hikoe Enomoto. Noboru Ito. Ryuzaburo Noda.
Andrew Bremner. Et al. Witt 4-(23, 7, 1). Witt 4-(23, 16, 52).

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Intersection numbers

For a tight $2s$ -(v, k, λ) design, zeros of the following polynomial are **intersection numbers**.

$$\Phi_s(x) = \sum_{i=0}^s (-1)^i \frac{\binom{v-s}{i} \binom{k-i}{s-i} \binom{k-i-1}{s-i}}{\binom{s}{i}} \binom{x}{i}.$$

Hermite polynomials

$$H_0(x) = 1, H_1(x) = x, \text{ and } H_{n+1}(x) = xH_n(x) - nH_{n-1}(x)$$

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- ▶ $s = 4$. Eiichi Bannai. Finitely many.

A thirteen degree polynomial

$$\begin{aligned} f(k, v) = & -3408102864 + 1506333312k^2 + 974873344k^4 - 488998144k^6 + 62323584k^8 - 3309568k^{10} + 65536k^{12} + 9310949028v - 1506333312kv - \\ & 4733985888k^2v - 1949746688k^3v - 1015706784k^4v + 1466994432k^5v + 511604992k^6v - 249294336k^7v - 49810560k^8v + 16547840k^9v + \\ & 1744896k^{10}v - 393216k^{11}v - 16384k^{12}v - 11097146016v^2 + 4733985888kv^2 + 6922441360k^2v^2 + 2031413568k^3v^2 - 1428764528k^4v^2 - \\ & 1534814976k^5v^2 + 209662720k^6v^2 + 199242240k^7v^2 - 21567744k^8v^2 - 8724480k^9v^2 + 786432k^{10}v^2 + 98304k^{11}v^2 + 7281931941v^3 - \\ & 5947568016kv^3 - 4944873072k^2v^3 + 412538336k^3v^3 + 1856597696k^4v^3 + 243542016k^5v^3 - 293538048k^6v^3 - 13016064k^7v^3 + \\ & 17194752k^8v^3 - 327680k^9v^3 - 253952k^{10}v^3 - 2755473732v^4 + 3929166288kv^4 + 1497511456k^2v^4 - 1155170432k^3v^4 - 582955856k^4v^4 + \\ & 183266304k^5v^4 + 58253568k^6v^4 - 16432128k^7v^4 - 1102464k^8v^4 + 368640k^9v^4 + 544096980v^5 - 1459281552kv^5 + 28759472k^2v^5 + \\ & 469164960k^3v^5 - 7038496k^4v^5 - 59703552k^5v^5 + 6536960k^6v^5 + 2050560k^7v^5 - 328320k^8v^5 - 18769932v^6 + 293023248kv^6 - \\ & 127930016k^2v^6 - 58917568k^3v^6 + 27050224k^4v^6 + 1258752k^5v^6 - 1642240k^6v^6 + 182784k^7v^6 - 14780538v^7 + 24513072kv^7 + \\ & 27560816k^2v^7 - 2875616k^3v^7 - 2296192k^4v^7 + 698880k^5v^7 - 61184k^6v^7 + 2961396v^8 - 764688kv^8 - 1582560k^2v^8 + 772608k^3v^8 - \\ & 143664k^4v^8 + 10752k^5v^8 - 191952v^9 + 203472kv^9 - 52816k^2v^9 + 7520k^3v^9 - 640k^4v^9 + 972v^{10} - 2352kv^{10} + 336k^2v^{10} + 45v^{11} \end{aligned}$$

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- ▶ $5 \leq s \leq 9$. Peter Dukes, Jesse Short-Gershman. None.

$$\beta = \frac{(v - k - s)\sqrt{(k - s + 1)(k - s)}}{(v - 2s + 1)^{3/2}}$$

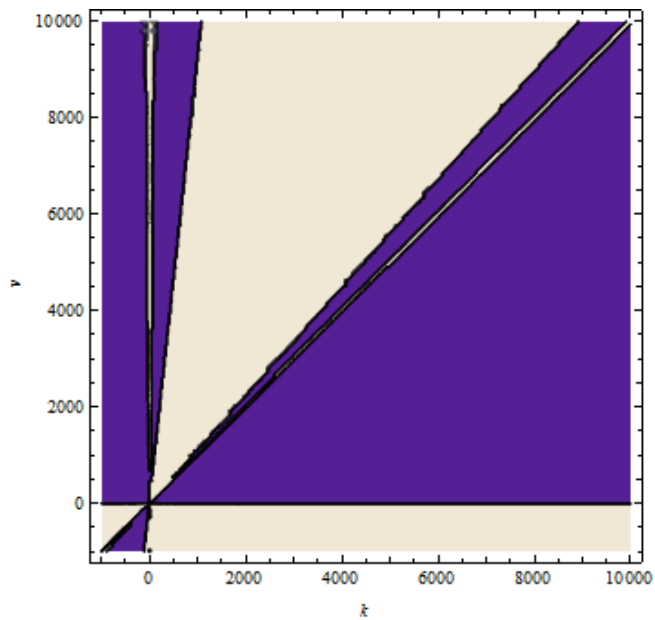
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β_0

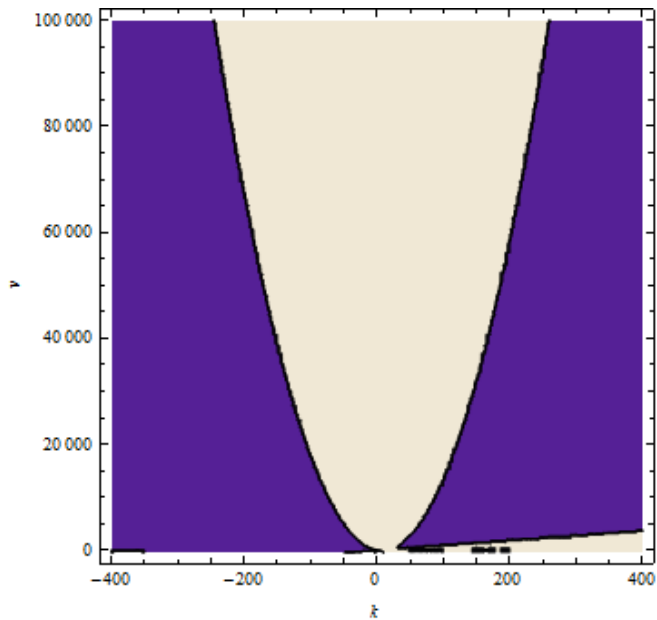
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Figure



Figure



Zeros

$$v = \frac{2}{1 - \sqrt[4]{\frac{3}{8}}}k + \frac{23}{500} \left(249 + 86\sqrt{6} + \sqrt{171312 + 70918\sqrt{6}} \right) + o(1)$$

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$$g(v, k) \in \mathbf{Z}$$

$$\lim_{v, k \rightarrow \infty} g(v, k) = \frac{9}{100} \left(6522 + 2808\sqrt{6} - \sqrt{56993328 + 24204417\sqrt{6}} \right)$$

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- ▶ $s = 4$. Ziqing Xiang. None.
- ▶ $s \geq 10$. ?

Thanks for your attention.